

Math 414 Lecture 25

Integer Programming

In many problems the variables must range over integers: the number of people you must hire, the number of chairs you must rent, a price in cents, ...

DEFINITION.

- A *pure integer* programming problem is a linear programming problem in which all the variables are required to be integers or natural numbers.
- In a *mixed integer* programming problem some of the variables may be required to be integral. The optimal value z , the coefficients a, b, c and the slacks are not required to be integers, just the variables.
- In a *0-1 programming* problem the variables must be only 0 or 1. *0-1 variables* are also called *Boolean* or *decision* variables.

Let Z = the set of integers, $N = \{0, 1, 2, 3, \dots\}$ be the set of natural numbers. Hence instead of x unrestricted or $x \geq 0$ we may require $x \in Z, x \in N$ or $x \in \{0, 1\}$.

Write the following as pure, mixed or 0-1 integer programming problems.

KNAPSACK PROBLEM. Let W = max weight you can carry on a hike. You have n items with weights w_1, \dots, w_n . Unfortunately $W < w_1 + \dots + w_n$, so you can't take them all. Assign values v_1, \dots, v_n to the n items such that the more important ones get higher values. Build a model for the problem of deciding for each item whether to take it or not in a way that maximizes the total value v .

Solution.

Let $x_i = 1$ if you decide to take the i th item,
 $= 0$ if you do not.

$$\begin{aligned} \text{Max } v &= v_1 x_1 + v_2 x_2 + \dots + v_n x_n, \\ \text{with} \\ w_1 x_1 + w_2 x_2 + \dots + w_n x_n &\leq W \\ x_1, \dots, x_n &\in \{0, 1\} \end{aligned}$$

ASSIGNMENT PROBLEM. There are n workers (one task each) w_1, \dots, w_n and n tasks (one worker each) t_1, \dots, t_n . r_{ij} = the revenue worker w_i generates when doing task t_j . Build a model for deciding which worker should do which task in order to maximize the total revenue r .

Solution. Let $x_{ij} = \begin{cases} 1 & \text{if } w_i \text{ does task } t_j, \\ 0 & \text{otherwise.} \end{cases}$

$$\text{max } r = \sum_{i,j} r_{ij} x_{ij}$$

with

$$\text{For } i = 1 \dots n, \quad i: x_{i1} + x_{i2} + \dots + x_{in} = 1 \quad \sum_{j=1}^n x_{ij} = 1, \text{ or } \leq$$

$$\text{For } j = 1 \dots n, \quad j: x_{1j} + x_{2j} + \dots + x_{nj} = 1 \quad \sum_{i=1}^n x_{ij} = 1, \text{ or } \leq$$

$$x_{ij} \in \{0, 1\} \quad \text{There are } 2n \text{ constraints, not } 2 \text{ constraints.}$$

STOCK CUTTING PROBLEM. Pipe comes in 8' and 12' lengths.

You need 10 3' lengths and 7 4' lengths. How many 8' and 12' lengths should be bought to minimize waste?

Solution.

Cutting patterns for 8' lengths.

| pattern # | pattern | waste | #times pattern is used |
|-----------|---------|-------|------------------------|
| 1: | 3 | +5 | x_1 |
| 2: | 4 | +4 | x_2 |
| 3: | 3+3 | +2 | x_3 |
| 4: | 4+3 | +1 | x_4 |
| 5: | 4+4 | +0 | x_5 |

Cutting patterns for 12' lengths.

| pattern # | pattern | waste | |
|-----------|---------|-------|----------|
| 6: | 3+3+3 | +3 | x_6 |
| 7: | 3+3+3+3 | +0 | x_7 |
| 8: | 4+3+3 | +2 | x_8 |
| 9: | 4+4+3 | +1 | x_9 |
| 10: | 4+4+4 | +0 | x_{10} |

Let x_j be the number of times pattern j is used.

min $w =$ _____ where w = total waste.
 with

$$\text{(three ft.) } t: \text{_____} = 10$$

$$\text{(four ft.) } f: \text{_____} = 7$$

$$x_1, \dots, x_{10} \in \mathbb{N}$$

TRAVELING SALESMAN PROBLEM. A salesman has to visit each of n cities c_1, \dots, c_n . He starts and ends at c_1 . He must visit each city exactly once. Suppose d_{ij} is the distance between cities c_i and c_j .

Find the route which minimizes the total distance.

Solution. Given a route, let $x_{ij} = 1$ if c_i and c_j are successive cities along the route; $x_{ij} = 0$ otherwise. For $i = 1, \dots, n$, let u_i = the position of city c_i (1st, 2nd, 3rd, ...) along the route. Let D = the total distance.

If the route is $c_1 \rightarrow c_4 \rightarrow c_2 \rightarrow c_3 \rightarrow c_1$,
 then $x_{14} = x_{42} = x_{23} = x_{31} = 1$, all other $x_{ij} = 0$,
 $u_1 = 1, u_4 = 2, u_2 = 3, u_3 = 4$.

$$\text{min } D = \sum \{d_{ij} x_{ij}; i, j = 1 \dots n\}$$

with

$$i: \sum_{j=1}^n x_{ij} = 1 \quad \text{for } i = 1 \dots n$$

$$j: \sum_{i=1}^n x_{ij} = 1 \quad \text{for } j = 1 \dots n$$

$$i-j: u_i - u_j + n x_{ij} \leq n-1 \quad i, j \in \{1, \dots, n\}, j \neq 1$$

$$x_{ij} \in \{0, 1\}, u_i \in \mathbb{N}$$

The last condition rules out short cut paths with two or more loops.

Full loop example (o.k.)

$$c_1 \rightarrow c_4 \rightarrow c_2 \rightarrow c_3 \rightarrow c_1$$

Two loop example (not o.k.)

$$c_1 \rightarrow c_4 \rightarrow c_1,$$

$$c_2 \rightarrow c_3 \rightarrow c_2$$